

THE PROOF CHECKERS PACHECK AND PASTÈQUE FOR THE PRACTICAL ALGEBRAIC CALCULUS

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FMCAD 2020

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Online

Formal Verification using Computer Algebra

Renewed interest in recent years

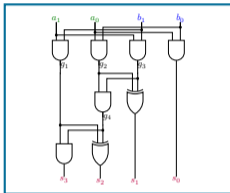
- A. Mahzoon, D. Große, and R. Drechsler. RevSCA: Using reverse engineering to bring light into backward rewriting for big and dirty multipliers. In DAC 2019.
- M. J. Ciesielski, T. Su, A. Yasin, and C. Yu. Understanding algebraic rewriting for arithmetic circuit verification: a bit-flow model. In IEEE TCAD 2019.
- A. Mahzoon, D. Große, C. Scholl, and R. Drechsler. Towards formal verification of optimized and industrial multipliers. In DATE 2020.

Algebraic reasoning in combination with SAT solving

- C. Bright, I. Kotsireas, and V. Ganesh. Applying computer algebra systems and SAT solvers to the Williamson conjecture. In Journal of Symbolic Computation, 2018.
- M. J. H. Heule. Computing small unit-distance graphs with chromatic number 5. CoRR, vol. abs/1805.12181, 2018.
- M. J. H. Heule, M. Kauers, and M. Seidl. Local search for fast matrix multiplication. In SAT 2019.
- D. Kaufmann, A. Biere, and M. Kauers. Verifying large multipliers by combining SAT and computer algebra. In FMCAD 2019.

Basic Idea of Algebraic Approach

System



Polynomials

$$B = \left\{ \begin{array}{l} x - a_0 * b_0, \\ y - a_1 * b_1, \\ s_0 - x * y, \\ \dots \end{array} \right\}$$

Specification

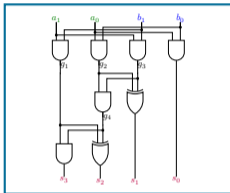
$$\sum_{i=0}^{2n-1} 2^i s_i - \left(\sum_{i=0}^{n-1} 2^i a_i \right) \left(\sum_{i=0}^{n-1} 2^i b_i \right)$$

Implication

$$\begin{array}{l} = 0 \quad \checkmark \\ \neq 0 \quad \times \end{array}$$

Basic Idea of Algebraic Approach

Multiplier



Polynomials

$$B = \left\{ \begin{array}{l} x - a_0 * b_0, \\ y - a_1 * b_1, \\ s_0 - x * y, \\ \dots \end{array} \right\}$$



Specification

$$\sum_{i=0}^{2n-1} 2^i s_i - \left(\sum_{i=0}^{n-1} 2^i a_i \right) \left(\sum_{i=0}^{n-1} 2^i b_i \right)$$

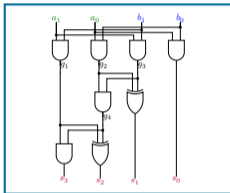


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Specification

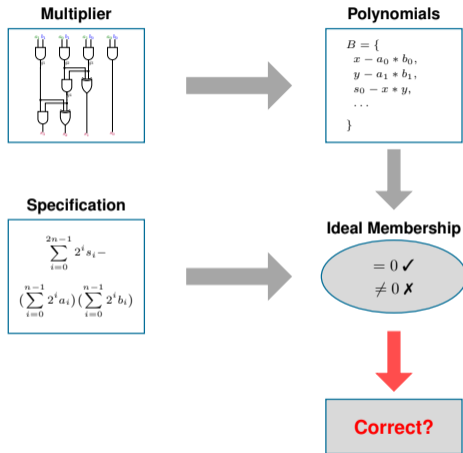
$$\sum_{i=0}^{2n-1} 2^i s_i - \left(\sum_{i=0}^{n-1} 2^i a_i \right) \left(\sum_{i=0}^{n-1} 2^i b_i \right)$$



Ideal Membership

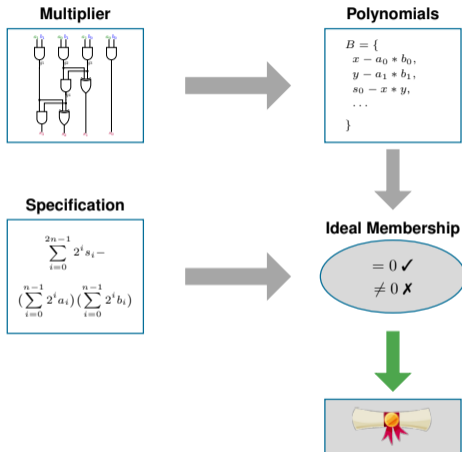
$$\begin{array}{l} = 0 \quad \checkmark \\ \neq 0 \quad \times \end{array}$$

Motivation



Problem: Verification might not be error free

Motivation



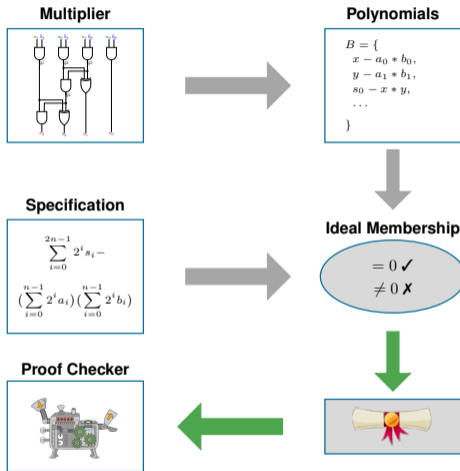
Problem: Verification might not be error free

Goal: Validate result of verification process

- Generate machine-checkable proofs
- Check by independent proof checkers

⇒ SC'2 2018: Practical Algebraic Calculus (PAC)
based on polynomial calculus

Motivation



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Contribution:

- Extending PAC:
indexing, deletion and extension rules
- Proof checker PACHECK and PASTÈQUE

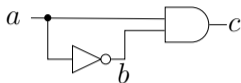
Ideal

Ideal. A nonempty subset $I \subseteq \mathbb{Z}[X]$ is called an ideal if

$$\forall p, q \in I : p + q \in I \quad \text{and} \quad \forall p \in \mathbb{Z}[X] \forall q \in I : pq \in I$$

Ideal membership problem. Given a polynomial $f \in \mathbb{Z}[X]$ and a (finite) set of polynomials $P \subseteq \mathbb{Z}[X]$, decide whether $f \in \langle P \rangle$, where $\langle P \rangle$ is the smallest ideal containing all elements of P , also known as the ideal *generated by* P .

Practical Algebraic Calculus - SC'2 2018

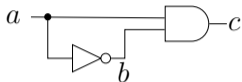


$$\begin{array}{ll}
 P = & -b+1-a, & b = \neg a \\
 & -c+a*b, & c = a \wedge b = a \wedge \neg a \\
 & -a^2+a & a = \perp \vee a = \top \\
 \text{Spec} = & c & c = \perp
 \end{array}$$

*	: -b+1-a,	a,	-a*b+a-a^2;
*	: -a^2+a,	-1,	a^2-a;
+	: -a*b+a-a^2,	a^2-a,	-a*b;
+	: -a*b,	-c+a*b,	-c;
*	: -c,	-1,	c;

$$\forall p, q \in I : p + q \in I \quad \text{and} \quad \forall p \in \mathbb{Z}[X] \forall q \in I : pq \in I$$

Practical Algebraic Calculus - SC'2 2018



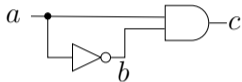
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$$\begin{array}{lll}
 * : & -b+1-a, & a, & -a*b+a-a^2; \\
 * : & -a^2+a, & -1, & a^2-a; \\
 + : & -a*b+a-a^2, & a^2-a, & -a*b; \\
 + : & -a*b, & -c+a*b, & -c; \\
 * : & -c, & -1, & c;
 \end{array}$$

Can be checked by our older proof checker PACTRIM.

1. Boolean Variables

Handle Boolean-value constraints implicitly to reduce number of proof steps.



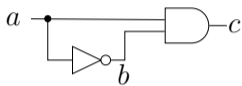
$$P = \begin{array}{l} -b+1-a, \\ -c+a*b, \\ \text{---}a^2+a \end{array}$$

$$\text{Spec} = c$$

$$\begin{array}{lll} * : & -b+1-a, & a, \quad -a*b; \\ + : & -a*b, & -c+a*b, \quad -c; \\ * : & -c, & -1, \quad c; \end{array}$$

2. Indices

Introduce indices to reduce proof size.



$$P = \begin{array}{l} 1 \quad -b+1-a; \\ 2 \quad -c+a*b; \end{array}$$

$$\text{Spec} = c$$

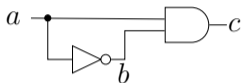
$$3 \quad * 1, \quad a, \quad -a*b;$$

$$4 \quad + 3, \quad 2, \quad -c;$$

$$5 \quad * 4, \quad -1, \quad c;$$

3. Deletion Rule

Introduce a deletion rule to reduce the memory usage of the proof checker.



$$P = \begin{array}{l} 1 \quad -b+1-a; \\ 2 \quad -c+a*b; \end{array}$$

$$\text{Spec} = c$$

$$3 \quad * 1, \quad a, \quad -a*b;$$

$$1 \quad d;$$

$$4 \quad + 3, \quad 2, \quad -c;$$

$$2 \quad d;$$

$$3 \quad d;$$

$$5 \quad * 4, \quad -1, \quad c;$$

4. Extension Rule

The extension rule allows to add model preserving polynomials to the constraint set.

$$\frac{\bar{x} \vee \bar{y} \quad y \vee z}{\bar{x} \vee z}$$

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$$\frac{xy \quad (1 - y)(1 - z)}{x(1 - z)}$$

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$$\frac{xy \quad (1-y)(1-z)}{x(1-z)}$$

$$\begin{aligned} P &= \begin{array}{l} 1 \quad x*y; \\ 2 \quad y*z-y-z+1; \end{array} \\ \text{Spec} &= \quad -x*z+x \end{aligned}$$

$$3 \quad = \quad f, \quad -z+1;$$

$$4 \quad * \quad 3, \quad y-1, \quad -f*y+f-y*z+y+z-1;$$

$$5 \quad + \quad 2, \quad 4, \quad -f*y+f;$$

EXT(i, v, p) $(X, P) \Rightarrow (X \cup \{v\}, P(i \mapsto -v + p))$
provided that $P(i) = \perp$ and $v \notin X$ and $p \in \mathbb{Z}[X]/\langle B(X) \rangle$,
and $p^2 - p \equiv 0 \pmod{\langle B(X) \rangle}$.

4. Extension Rule

The extension rule allows to add model preserving polynomials to the constraint set.

$$\frac{xy \quad (1-y)(1-z)}{x(1-z)}$$

$$P = \begin{array}{l} 1 \quad x*y; \\ 2 \quad y*z-y-z+1; \end{array}$$

$$\text{Spec} = -x*z+x$$

$$\begin{array}{l} 3 = f, -z+1; \\ 4 * 3, y-1, -f*y+f-y*z+y+z-1; \\ 5 + 2, 4, -f*y+f; \\ 6 * 1, f, f*x*y; \\ 7 * 5, x, -f*x*y+f*x; \\ 8 + 6, 7, f*x; \\ 9 * 3, x, -f*x-x*z+x; \\ 10 + 8, 9, -x*z+x; \end{array}$$

PACHECK

- 1700 lines of C code
- supports new and old PAC format
- PACHECK reads three input files `<input>`, `<proof>`, and `<target>`.
- Verifies that the polynomial in `<target>` is contained in the ideal generated by the polynomials in `<input>` using the rules provided in `<proof>`.

PACHECK

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- supports new and old PAC format
- PACHECK reads three input files `<input>`, `<proof>`, and `<target>`.
- Verifies that the polynomial in `<target>` is contained in the ideal generated by the polynomials in `<input>` using the rules provided in `<proof>`.
- Polynomial arithmetic is implemented from scratch.
- A polynomial is represented as a linked list of monomials:
 - Coefficients are represented using the GMP library.
 - Terms are ordered linked lists of variables.

Variable ordering

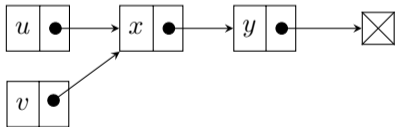
- The order of the variables has an enormous effect on the memory usage:
 - Terms are internally shared

Variable ordering

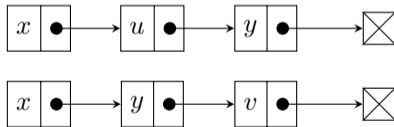
- The order of the variables has an enormous effect on the memory usage:
 - Terms are internally shared

Assume we want to represent the terms uxy and vxy .

$$v > u > x > y$$



$$x > u > y > v$$



Variable ordering

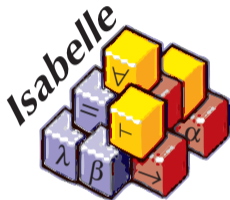
- The order of the variables has an enormous effect on the memory usage:
 - Terms are internally shared
- Example: 7 mio. rules, 50% memory increase using a different ordering

Implemented Orderings:

- Default: Variables are ordered using `strcmp`.
- Alternative: Same variable ordering as used in input files.
- Both orderings can also be reversed.

PASTÈQUE

Theorem Prover Isabelle/HOL



Refinement Approach, relying on Isabelle's Refinement Framework

- abstract specification on ideals: specification in ideal
- final step: executable checker

Isabelle's Archive of Formal Proofs 8 000 lines of code

Refinement

$$\text{no error} \implies \text{spec} \in \langle P \rangle \wedge \langle P \rangle|_X \subseteq \langle P_0 \rangle|_X$$

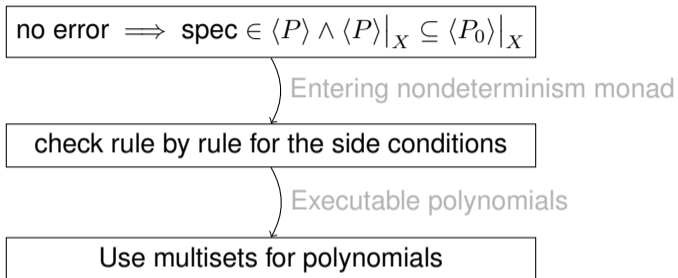
Refinement

no error $\implies \text{spec} \in \langle P \rangle \wedge \langle P \rangle|_X \subseteq \langle P_0 \rangle|_X$

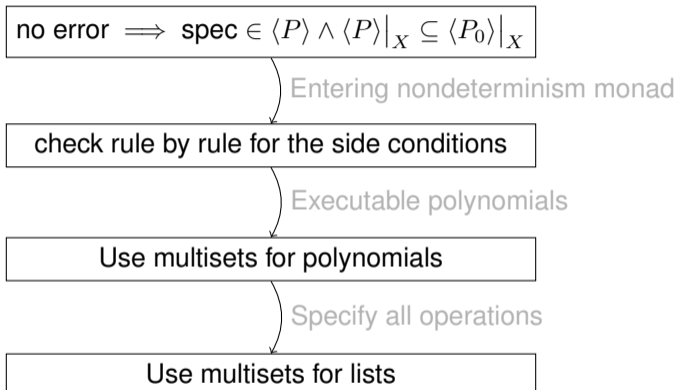
) Entering nondeterminism monad

check rule by rule for the side conditions

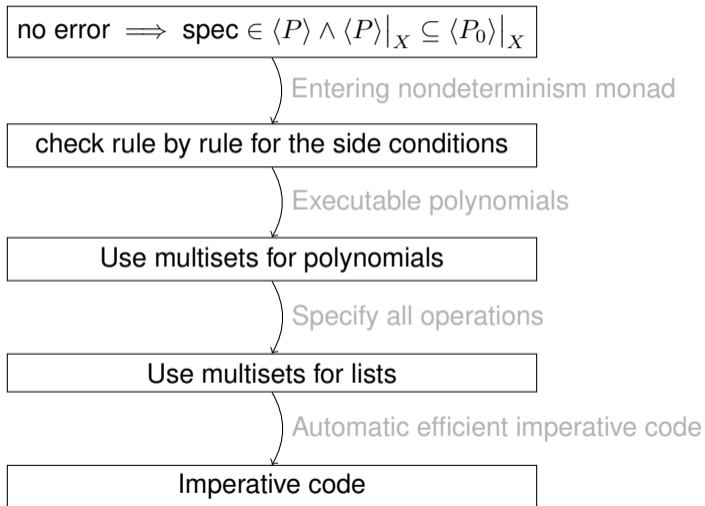
Refinement



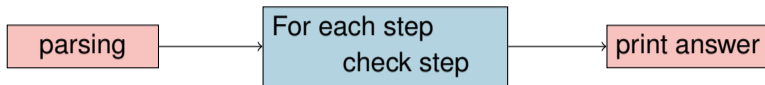
Refinement



Refinement

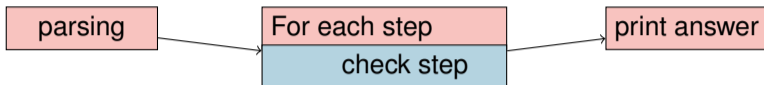


Refinement



- Translation to Standard ML is trusted
- Parser and plumbing (including answer printing) also trusted

Parsing vs Loop



Alternative, `uloop` variant:

- same functions to check the steps
- but: hand-written (trusted) version of the loop that iterates through the steps

Advantage: more memory efficient

Tool Demonstration

$$\frac{xy(1-y)(1-z)}{x(1-z)}$$

$$P = \begin{array}{l} 1 \ x*y; \\ 2 \ y*z-y-z+1; \end{array}$$

$$\text{Spec} = -x*z+x$$

```
3 = f, -z+1;
4 * 3, y-1, -f*y+f-y*z+y+z-1;
5 + 2, 4, -f*y+f;
2 d;
4 d;
6 * 1, f, f*x*y;
1 d;
7 * 5, x, -f*x*y+f*x;
5 d;
8 + 6, 7, f*x;
6 d;
7 d;
9 * 3, x, -f*x-x*z+x;
3 d;
10 + 8, 9, -x*z+x;
```



```
$ pacheck ex1.{input,proof,target}
[pacheck] Pacheck Version 001
[pacheck] Practical Algebraic Calculus Proof Checker
[pacheck] Copyright (C) 2020, Daniela Kaufmann, Johannes Kepler University Linz
[pacheck] compressed mode with indices assumed
[pacheck] sorting according to strcmp
[pacheck] checking target enabled
[pacheck] reading target polynomial from 'ex1.target'
[pacheck] read 8 bytes from 'ex1.target'
[pacheck] reading original polynomials from 'ex1.input'
[pacheck] found 2 original polynomials in 'ex1.input'
[pacheck] read 20 bytes from 'ex1.input'
[pacheck] reading polynomial algebraic calculus proof from 'ex1.proof'
[pacheck] found and checked 8 inferences in 'ex1.proof'
[pacheck] read 219 bytes from 'ex1.proof'
[pacheck] found 1 target polynomial inference
[pacheck] proof length 10 (number of polynomials)
[pacheck] proof size 25 (on average 2.5 terms per polynomial)
[pacheck] proof degree 3 (internal maximum degree 3)
[pacheck] searched 32 inferences 0.1 average collisions
[pacheck] 10 inferences, 3.2 average searches
[pacheck] original inferences 2 (20% of total rules)
[pacheck] inference rules 8 (80% of total rules)
[pacheck] addition inference rules 3 (38% of inference rules)
[pacheck] multiplication inference rules 4 (50% of inference rules)
[pacheck] extension rules 1 (12% of inference rules)
[pacheck] deletion inference rules 3 (30% of total rules)
[pacheck] maximum 9 of total 10 terms (90%)
[pacheck] searched 52 terms 81% hits 0.3 average collisions
[pacheck] maximum 2261 bytes allocated (0.0 MB)
[pacheck] maximum resident set size 5922816 bytes (5.6 MB)
[pacheck] process time 0.001 seconds
[pacheck] TARGET CHECKED
```

```
$ pasteque ex1.{input,proof,target}
c polys parsed
c *****
c pac parsed
c spec parsed
c Now checking
s SUCCESSFULL
c
c
c ***** stats *****
c parsing polys file init (nonGC): 0.000 s = 0.000 s (usr) 0.000 s (sys)
c parsing pac file init (nonGC): 0.000 s = 0.000 s (usr) 0.000 s (sys)
c full init (nonGC): 0.000 s = 0.000 s (usr) 0.000 s (sys)
c time solving (nonGC): 0.000 s = 0.000 s (usr) 0.000 s (sys)
c time GC: 0.000 s = 0.000 s (usr) 0.000 s (sys)
c time solving(full): 0.000 s
c Overall (nonGC): 0.000 s = 0.000 s (usr) 0.000 s (sys)
c overall GC: 0.000 s = 0.000 s (usr) 0.000 s (sys)
c Overall(full): 0.000 s
```

```
$ vim ex1.proof
```

```
3 = f, -z+1;
4 * 3, y-1, -f*y+f-y*z+y+z-1;
5 + 2, 4, -f*y+f;
2 d;
4 d;

6 * 1, f, f*x * y;
1 d;

7 * 5, x, -f*x*y+f*x;
8 + 6, 7, f*x;
9 * 3, x, -f*x-x*z+x;
10 + 8, 9, -x*z+x;
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3 = f, -z+1;
4 * 3, y-1, -f*y+f-y*z+y+z-1;
5 + 2, 4, -f*y+f;
2 d;
4 d;

6 * 1, f, f*x + y;
1 d;

7 * 5, x, -f*x*y+f*x;
8 + 6, 7, f*x;
9 * 3, x, -f*x-x*z+x;
10 + 8, 9, -x*z+x;
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[pacheck] read 8 bytes from 'ex1.target'
[pacheck] reading original polynomials from 'ex1.input'
[pacheck] found 2 original polynomials in 'ex1.input'
[pacheck] read 20 bytes from 'ex1.input'
[pacheck] reading polynomial algebraic calculus proof from 'ex1.proof'
*** 'pacheck' error in polynomial multiplication rule 4 with index 6 in 'ex1.proof' line 7: conclusion polynomial:
f*x+y
does not match expected result:
f*x*y
```

Evaluation

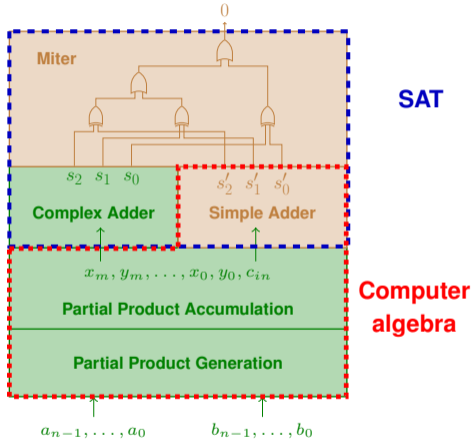
Comparison between:

PACTRIM the older version

PACHECK with no delete, no index, everything

PASTÈQUE with and without `uloop`

Evaluation



Complex adder: Replace it by simpler adder [DATE20, Kaufmann et al.] via extensions.

Evaluation Results

multiplier	n	steps (10^6)	ext	deg	PACTRIM		no index		PACHECK		no delete		default	
					sec	MB	sec	MB	sec	MB	sec	MB	sec	MB
					btor	128	0.4	0	3	10	105	11	100	5
btor	256	1.6	0	3	60	459	62	435	25	1 144	25	364		
btor	512	6.3	0	3	395	2 066	402	1 972	138	4 956	141	1 461		
sp-ar-rc	128	0.6	0	4	16	156	16	148	6	454	6	136		
sp-ar-rc	256	2.3	0	4	92	687	96	651	29	1 858	27	541		
sp-ar-rc	512	9.4	0	4	587	3 107	617	2 965	146	7 683	134	2 171		
sp-ar-cl	32	1.6	2 081	256	31	405	36	354	23	773	21	353		
sp-dt-lf	32	0.3	1 684	46	3	82	3	73	2	122	2	73		
bp-ct-bk	32	0.2	1 405	25	2	57	2	52	1	86	1	51		
bp-wt-cl	32	5.6	5 740	764	242	1 716	302	1 430	193	4 324	181	1 428		

Evaluation Results

multiplier	n	steps (10^6)	ext	deg	PACHECK		PASTÈQUE			
					default		default		uloop	
					sec	MB	sec	MB	sec	MB
btor	128	0.4	0	3	5	92	22	3 886	17	1 773
btor	256	1.6	0	3	25	364	105	21 157	79	4 364
btor	512	6.3	0	3	141	1 461	531	64 412	416	22 292
sp-ar-rc	128	0.6	0	4	6	136	31	5 002	23	1 608
sp-ar-rc	256	2.3	0	4	27	541	139	32 525	102	8 769
sp-ar-rc	512	9.4	0	4	134	2 171	608	64 412	471	25 632
sp-ar-cl	32	1.6	2 081	256	21	353	121	40 654	113	9 492
sp-dt-lf	32	0.3	1 684	46	2	73	11	1 679	11	886
bp-ct-bk	32	0.2	1 405	25	1	51	8	1 600	7	1 068
bp-wt-cl	32	5.6	5 740	764	181	1 428	786	58 867	774	64 404

Conclusion & Future Work

■ Proof Checker:

- PACHECK is 30-80% faster than older proof checker.
- PASTÈQUE is 4×-slower, but fully verified.

■ New PAC format:

- No need to add extension variables to input polynomials anymore.
- Indices and assumption of Boolean variables reduces the proof size.
- Deletion rules reduce the memory usage.

■ Investigate whether lifting the restrictions on EXT is useful.

■ Explore connection to Nullstellensatz Proofs to gain shorter proofs:

- D. Kaufmann and A. Biere. Nullstellensatz-Proofs for Multiplier Verification. In CASC 2020.

THE PROOF CHECKERS PACHECK AND PASTÈQUE FOR THE PRACTICAL ALGEBRAIC CALCULUS

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